#### **Admin**

Lab | => Assign | We're off and running!



# Today: From Assembly to C (and back again)

Branch instructions, PC-relative

C language as "high-level" assembly

What does a compiler do?

**Makefiles** 

## Control flow, pc register

Instructions stored in contiguous memory

pc tracks address in memory where instructions are being read

pc register separate from x0-x31, not accessible to most instructions, use special instructions to access/change pc

Default is "straight-line" code: next instruction to execute is at next higher memory address (pc = pc + 4)

jump instruction assigns pc to different address

j target

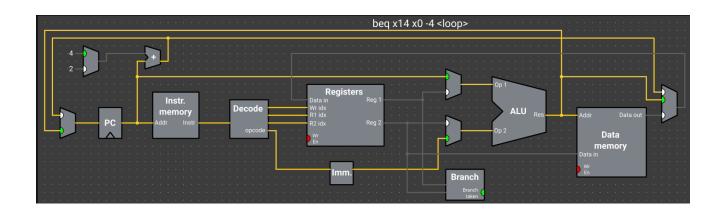
Jump is unconditional (always taken)

Branch is conditionally taken based on test

#### **Branch instructions**

Mnemonic	Action
BEQ rs1,rs2,imm12	Branch equal (rs1 = rs2)
BNE rs1,rs2,imm12	Branch not equal (rs1 ≠ rs2)
BGE rs1,rs2,imm12	Branch greater than or equal (rs1 ≥ rs2)
BLT rs1,rs2,imm12	Branch less than (rs1 < rs2)

Q: How to... branch greater? Branch less-equal? Branch zero? Branch negative?



If condition satisfied, branch is taken (pc = pc + imm12) otherwise falls through (pc = pc + 4)

## Challenge for you

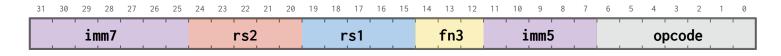
# Write an assembly program to count the "on" bits in a given numeric value

```
li a0, val
li a1, 0

// a0 holds input value
// use a1 to store count of on bits in value
```

```
1 ms
                             $ > /·
        Source code Input type: ● Assembly ○ □ Executable code View mode: ○ Binary ● Disassembled
         1 li a0.0x54
                                                                                 addi x10 x0 84
         2 li a1.0
                                                                                 addi x11 x0 0
         3 more:
              andi a2,a0,1
                                                       00000000000000008 <more>:
              add a1,a1,a2
                                                                                 andi x12 x10 1
              srli a0.a0.1
                                                                                 add x11 x11 x12
              bne a0, zero, more
                                                                    00155513
                                                                                  srli x10 x10 1
                                                                    fe051ae3
                                                                                  bne x10 x0 -12 <more>
```

## Branch instruction encoding



if rs1 cmp\_op rs2 pc = pc + imm12

- branch target computed as PC-relative offset
- purple bits encode offset (immediate)
- "position-independent" code

I2-bit immediate expressed as count of 2-byte steps Q: How far can this reach?

## ISA design is an art form!

As much about what is **omitted** as what is **included** 

All registers general-purpose registers, no act on memory ("load-store")

Simplicity (avoid redundancies, single addressing mode)

Isolate architecture from implementation (no delay slots branch/load, no condition codes)

Regularity: all instructions 4-bytes (2 for compressed)

Handling/placement of bits in encoding for ease of decode/data path

Modular, extensible (tiny base ISA, orthogonal additions)

Data-informed design (learn from past)

#### Why assembly?

What you see is what you get

No surprises

Precise control, timing

Unfettered access to hardware

**But**... tedious, hard to read, hardware-specific

### Why C?

More concise

Easier to read

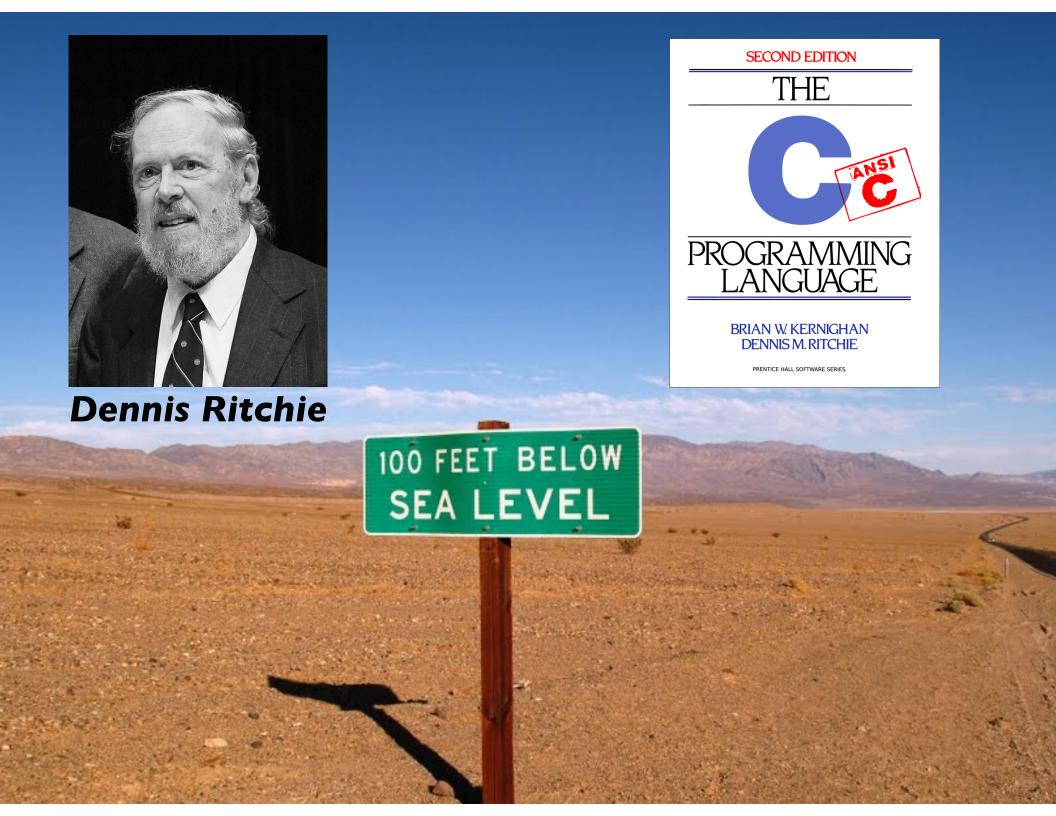
Can name variables and structures

Type-checking

More portable, architecture-neutral

**Functions** 

Real question is not whether to use assembly, but when...



# C is the language of choice for systems programmers



Ken Thompson built UNIX using C

This is not coincidence!

C features closely model the ISA: data types, arithmetic/logical operators, control flow, access to memory, ... all provided in form of portable abstractions

"BCPL, B, and C family of languages are particularly oriented towards system programming, are small and compactly described, and are amenable to translation by simple compilers. They are "close to the machine" in that the abstractions they introduce are readily grounded in the concrete data types and operations supplied by conventional computers, and they rely on library routines for input-output and other interactions with an operating system. . . . At the same time, their abstractions lie at a sufficiently high level that, with care, portability between machines can be achieved."

— Dennis Ritchie

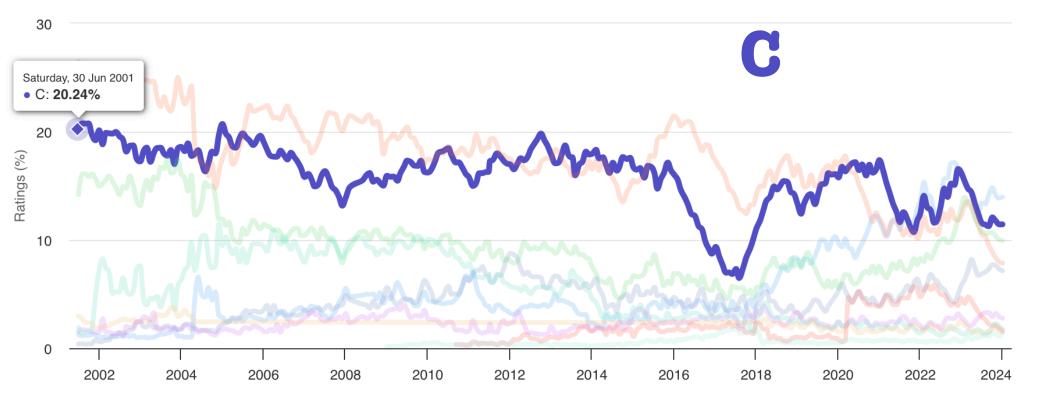
## The C Programming Language

- "C is quirky, flawed, and an enormous success"
- Dennis Ritchie
- "C gives the programmer what the programmer wants; few restrictions, few complaints"
- Herbert Schildt
- "C: A language that combines all the elegance and power of assembly language with all the readability and maintainability of assembly language"
- Unknown

## Programming language popularity over time

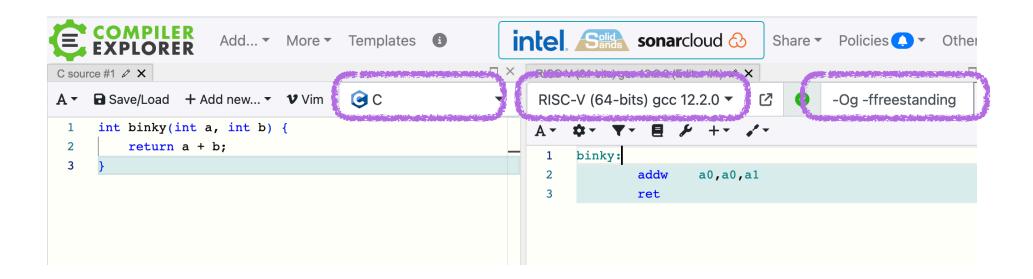
#### **TIOBE Programming Community Index**

Source: www.tiobe.com



#### **Compiler Explorer**

is a neat interactive tool to see translation from C to assembly. Let's try it now!



https://godbolt.org

```
Configure settings to follow along:
C
RISC-V (64 bits) gcc 12
-Og -ffreestanding
```

## Major props to the C compiler

#### Higher-level abstractions, structured programming

Named variables, constants

Arithmetic/logical operators

Control flow

#### **Portable**

Not tied to particular ISA or architecture

#### Low-level enough to get to machine when needed

Bitwise operations

Direct access to memory

Embedded assembly, too!

## Compile-time vs. runtime

Compile-time: compiler running on your laptop

- read C source text, parse/check semantically valid
- analyze code to understand structure/intent
- generate assembly instructions, assembler to binary

Runtime: program binary running on Pi

- load machine instructions to memory
- fetch/decode/execute

Optimizer does work at CT to streamline count of instructions to be executed at RT

## Know your tools!

#### **Assembler** as

```
Transform assembly code (text)
into object code (binary machine instructions)
Mechanical rewrite, few surprises
```

```
Compiler gcc
Transform C code (text)
into object code
(likely staged C → asm → object)
Complex translation, high artistry
```

### **Make**

One-step build process using make

Makefile is text file that describes build steps as "recipes"

Dependencies determine which steps needed to re-build

```
Rule

blink.bin: blink.s

riscv64-unknown-elf-as blink.s -o blink.o

riscv64-unknown-elf-objcopy blink.o -0 binary blink.bin

Target

run: blink.bin

Dependency

mango-run blink.bin
```

Writing out explicit recipes becomes onerous, so make has all kinds of ways to match patterns, define variables, etc.

## Make pattern rules

```
NAME = myprogram
ARCH = -march=rv64imac -mabi=lp64
CFLAGS = $(ARCH) -g -Og -Wall -ffreestanding
LDFLAGS = $(ARCH) -nostdlib
all: $(NAME).bin
%.bin: %.elf
        riscv64-unknown-elf-objcopy $< -0 binary $@</pre>
%.elf: %.o
        riscv64-unknown-elf-gcc $(LDFLAGS) $< -o $@</pre>
%.o: %.c
        riscv64-unknown-elf-gcc $(CFLAGS) -c $< -o $@
```

### Bare-metal vs. Hosted

The default build process for C assumes a hosted environment.

What does a hosted system have that we don't?

- standard libraries
- standard start-up sequence
- OS services

To build bare-metal, our Makefile disables these defaults We supply our own replacements where needed

## **Build settings for bare-metal**

Compile freestanding

```
CFLAGS = -ffreestanding
```

Link excludes standard library and start files

```
LDFLAGS = -nostdlib
```

Link with gcc if need software floating point

```
LDLIBS = -lgcc
```

Write our own code for all libs and start files

This puts us in an exclusive club...

```
int getRandomNumber()
{
    return 4; // chosen by fair dice roll.
    // guaranteed to be random.
}
```